

Qubits Europe 2019

# Flight Gate Assignment with a Quantum Annealer

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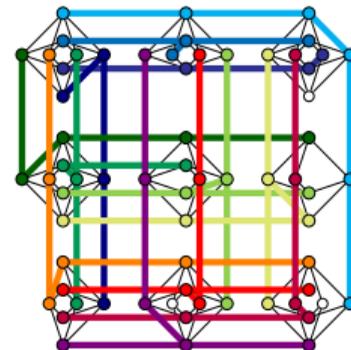
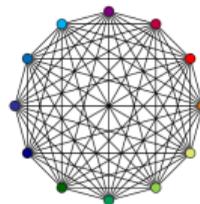
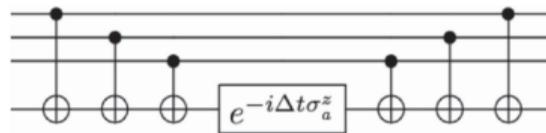
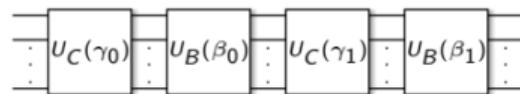
26th March 2019

A large, curved view of the Earth from space, showing the blue atmosphere, white clouds, and green and brown landmasses. The view is centered on Europe and the Mediterranean region.

Knowledge for Tomorrow

# Algorithmic Quantum Computing Research at DLR

- Quantum Optimization Algorithms
- Quantum Compiling
- Embedding strategies for Quantum Annealing
  - Complete graph in broken Chimera
  - Weight distribution problem



# Aerospace Applications at DLR

for Quantum Annealing

- Air Traffic Management
- Satellite Telemetry Verification
- Earth Observation Mission Planning
- Flight Gate Assignment

for Gate-Based Quantum Computing

- QAOA for scheduling problems
- HHL for Radar Cross Section
- Quantum Simulation for Battery Research



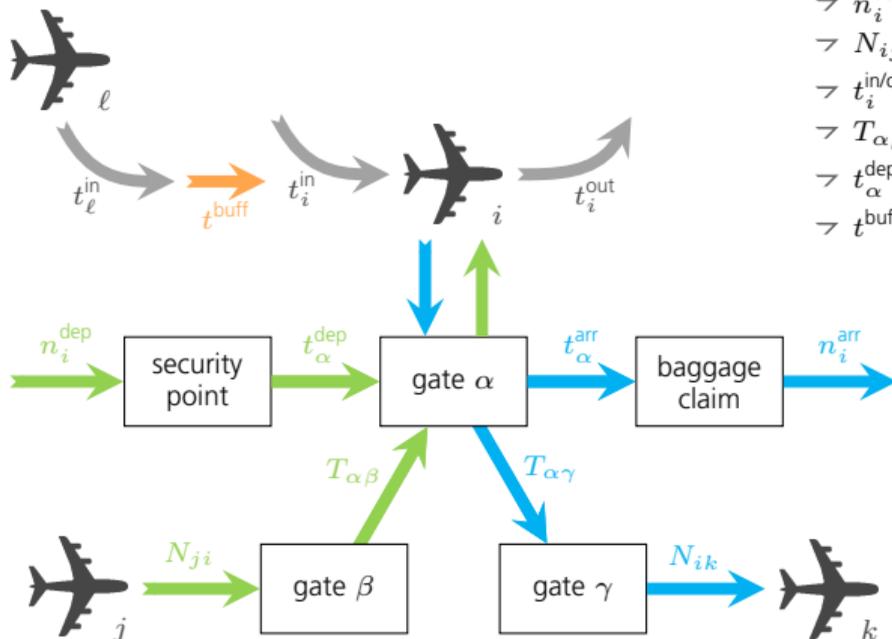
# Flight Gate Assignment

## A day at Frankfurt Airport

- about 1300 aircraft movements (arrival and departure)
- more than 90% are passenger flights
- more than 170000 passengers
- about 60% transfer passengers
- 278 gates



# Passenger Flows



- $F, G$  sets of flights and gates
- $n_i^{\text{dep/arr}}$  passengers which depart/ arrive with flight  $i$
- $N_{ij}$  transfer passengers from flight  $i$  to  $j$
- $t_i^{\text{in/out}}$  arrival/departure time of flight  $i$
- $T_{\alpha\beta}$  average time to get from gate  $\alpha$  to  $\beta$
- $t_\alpha^{\text{dep/arr}}$  average time to arrive at/ leave from gate  $\alpha$
- $t^{\text{buff}}$  buffer time between two flights at the same gate

Which flight should be assigned to which gate, such that the total transit time of the passengers is minimal?

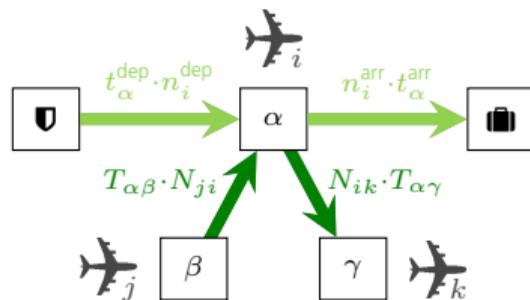
$$A : F \rightarrow G$$



# FGA Binary Program

Variables  $x \in \{0, 1\}^{F \times G}$  with

$$x_{i\alpha} = \begin{cases} 1, & \text{if flight } i \text{ takes gate } \alpha, \\ 0, & \text{otherwise} \end{cases}$$



Minimizing the total transfer time with objective function

$$\begin{aligned} T(x) &= T_{\text{arr}}(x) && + T_{\text{dep}}(x) && + T_{\text{transfer}}(x) \\ &= \underbrace{\sum_{i\alpha} n_i^{\text{arr}} t_{\alpha}^{\text{arr}} x_{i\alpha} + \sum_{i\alpha} n_i^{\text{dep}} t_{\alpha}^{\text{dep}} x_{i\alpha}}_{\text{linear}} + \underbrace{\sum_{ij\alpha\beta} N_{ij} T_{\alpha\beta} x_{i\alpha} x_{j\beta}}_{\text{quadratic}} \end{aligned}$$

⇒ Quadratic Assignment Problem

- fundamental problem in combinatorial optimization, NP-hard
- seems to exploit possible advantages of the D-Wave machine



# Constraints and Penalty Terms

1. One gate per flight

$$\sum_{\alpha} x_{i\alpha} = 1 \quad \forall i \in F$$

2. Different gates if standing times of two flights overlap forbidden pairs

$$P = \left\{ (i, j) \in F^2 : t_i^{\text{in}} < t_j^{\text{in}} < t_i^{\text{out}} + t^{\text{buff}} \right\}$$

$$x_{i\alpha} + x_{j\alpha} \leq 1 \Leftrightarrow x_{i\alpha} \cdot x_{j\alpha} = 0 \quad \forall (i, j) \in P \quad \forall \alpha \in G$$

$$\Rightarrow \text{Penalty terms } C_{\text{one}}(x) = \sum_i \left( \sum_{\alpha} x_{i\alpha} - 1 \right)^2,$$

$$C_{\text{not}}(x) = \sum_{\alpha} \sum_{(i,j) \in P} x_{i\alpha} x_{j\alpha} \quad \text{where } C_{\text{one/not}} \begin{cases} > 0, & \text{if constraint is violated} \\ = 0, & \text{if constraint is fulfilled} \end{cases}$$



# QUBO with Penalty Weights

$$Q(x) = T(x) + \lambda_{\text{one}} C_{\text{one}}(x) + \lambda_{\text{not}} C_{\text{not}}(x)$$

Need to ensure that a solution always fulfills constraints, hence  $\Delta C > \Delta T$

⇒ Comparing coefficients in worst cases for

→ not assigning a flight to any gate

$$\lambda_{\text{one}} > \max_{i,\alpha} \left( n_i^{\text{dep}} t_{\alpha}^{\text{dep}} + n_i^{\text{arr}} t_{\alpha}^{\text{arr}} + \max_{\beta} T_{\alpha\beta} \sum_j N_{ij} \right)$$

→ assigning a pair of forbidden flights to the same gate

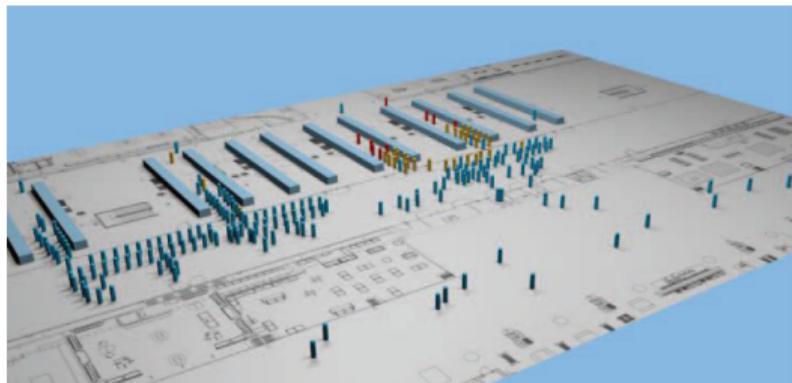
$$\lambda_{\text{not}} > \max_{i,\alpha,\gamma} \left( n_i^{\text{dep}} t_{\alpha}^{\text{dep}} - n_i^{\text{dep}} t_{\gamma}^{\text{dep}} + n_i^{\text{arr}} t_{\alpha}^{\text{arr}} - n_i^{\text{arr}} t_{\gamma}^{\text{arr}} + \max_{\beta} (T_{\alpha\beta} - T_{\gamma\beta}) \sum_j N_{ij} \right)$$

⇒ Refinement by bisection of weights yielding valid or invalid solutions



# Airport Data

- Flight schedule for one day from a mid-sized European airport
- Passenger flow from agent-based simulation of Martin Jung
- Extracted total instance: 293 flights and 97 gates
- ⇒ Over 28000 binary variables with about 400 Mio. couplings



M. Jung et al. (DLR-FW)

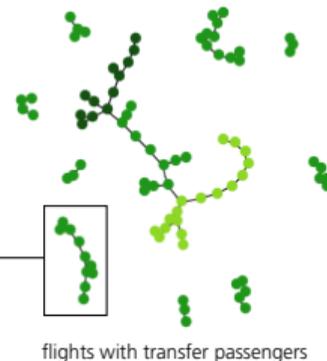
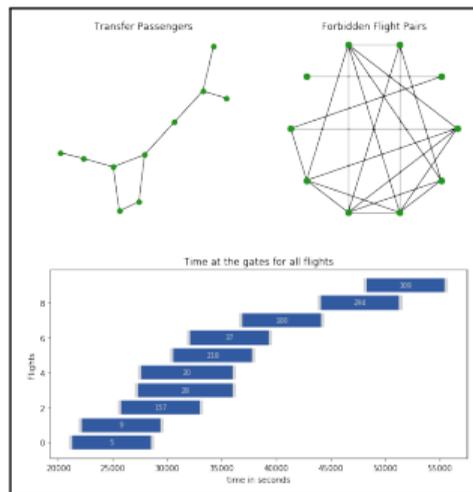
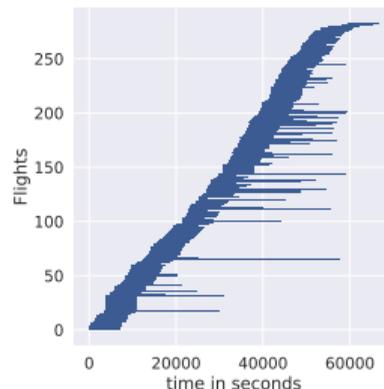


# Instance Preprocessing

- Splitting too long on-block times
- Reducing to only flights with transfers
- Extracting connected subgraphs
- Further slicing of largest subgraph randomly

⇒ 163 instances:

- 3 to 16 flights
- 2 to 16 gates



# Bin Packing

➤ Maximum coefficient ratio of QUBO  $C_Q = \frac{\max_{ij} |Q_{ij}|}{\min_{ij} |Q_{ij}|}$

➤ Reducing maximum coefficient ratio to overcome precision issues

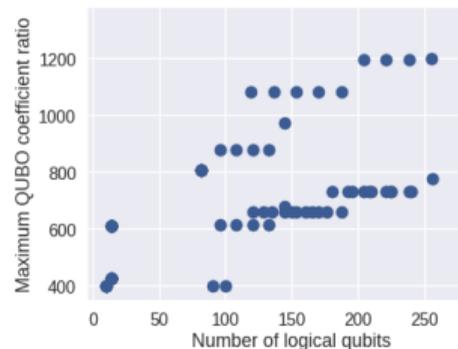
$$T_{\alpha\beta}, t_{\alpha}^{arr}, t_{\alpha}^{dep} \rightarrow \{0, 1, \dots, T\} \text{ with } T \in \{2, 3, 6, 10\}$$

$$N_{ij}, n_i^{arr}, n_i^{dep} \rightarrow \{0, 1, \dots, N\} \text{ with } N \in \{2, 3, 6, 10\}$$

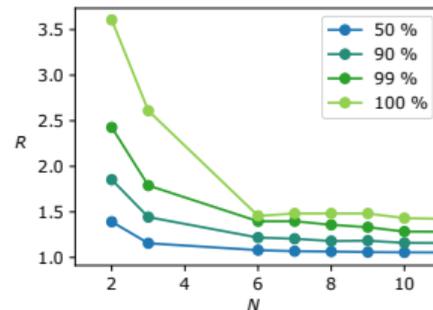
➤ Approximation ratio (solved with SCIP)

$$R = \frac{Q(\arg \min_x \bar{Q}(x))}{\min_x Q(x)}$$

⇒ Little effect on solution quality



$$\Rightarrow C_{\bar{Q}} \ll C_Q$$



# Annealing Setup

## ➤ Embedding

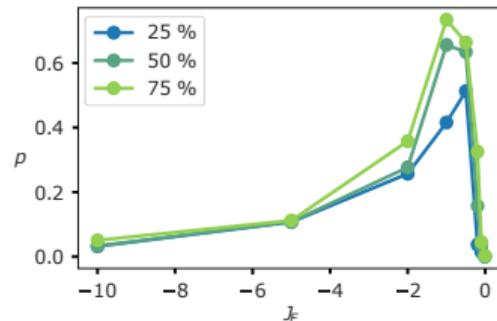
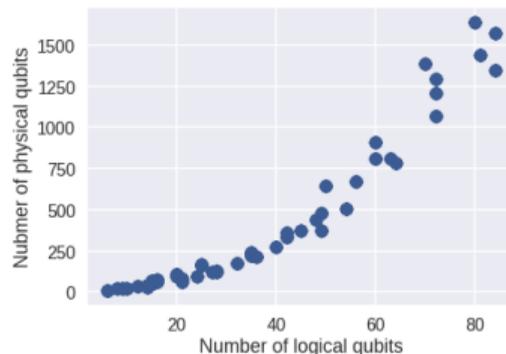
- Quadratic overhead
- Up to 84 logical qubits  
(# Variables = #Flights · #Gates)

## ➤ Intra-logical coupling ( $J_F$ )

- Influences success probability  $p$
- Best option by scanning: -1 in units of largest coefficient

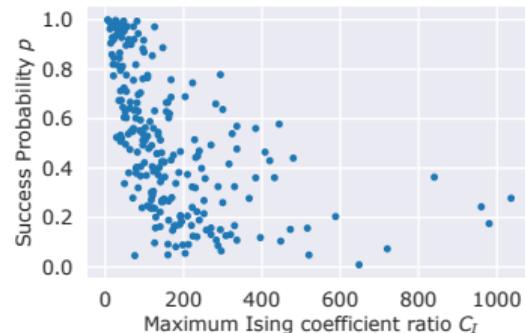
## ➤ (Standard) Run parameters

- Annealing time  $20\mu s$  with 1000 runs
- Majority voting



# Annealing Results

- QUBO to Ising transformation
  - increases maximum coefficient ratio significantly
  - ⇒ large Ising coefficients suppress success probability
  
- Time to solution with 99% certainty  $T_{99} = \log_{1-p}(1 - 0.99)T_{\text{anneal}}$ 
  - grows with problem size → because of larger coefficients?
  - due to small problem sizes asymptotic behaviour unclear



# Summary

- Flight gate assignment is amenable to QA
- Precision issues due to large coefficients
- Mitigate limited precision by bin packing
- Open questions:
  - How to recombine partial solutions?
  - How would larger instance perform?
  - Are these instances hard for classical solvers?



## Questions?

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## Related Article:

*Flight Gate Assignment with a Quantum Annealer*

T. Stollenwerk, E. Lobe and M. Jung, QTOP, Springer, 2019



Knowledge for Tomorrow



DLR